## Math 466/666 Programming Project 1

1. List the members on your team and explain how the work for this project was conducted. Please provide details concerning how many meetings were held, what was discussed at each meeting and what work was done between meetings. Further include a written statement attesting that the submitted report represents the original efforts of the team members listed and, in particular, does contain the work of other students in the class.

You answer will be different than the made up one here.

This project was completed in four meetings by Student A, Student B and Student C. Setting up a discussion group on Web Campus helped schedule these meetings; however, email was used to send working copies of the project report back and forth between members of the team.

The initial meeting was held through Zoom. Student C had trouble connecting and then their computer crashed so the meeting was abruptly cancelled and rescheduled for the same time on the following week.

At the second meeting Test Student A wrote code to calculate the recurrence relation in question 2 using the terminal version of Julia as JupyterLAB was not installed. Code for this relation was communicated through the screen share and later emailed to each team member to work with. Student B agreed to watch the video on installing JupyterLAB. Student C volunteered to perform a web search to find additional information on the Foias constant.

At the third meeting most of the time was spent helping Student B install Julia on their computer using a screen share over Zoom. While trying to write a loop to determine whether the sequence was bounded, Anaconda started using all available computer cycles and then Zoom crashed. It was decided to set up the another meeting through Web Campus.

At the fourth meeting Student C presented two interesting links which discussed the Foias constant. A draft was started for the project report using a Jupyter LAB notebook by Student C and sent through email to each member of the team.

At the fifth meeting, new code was written and improved by all team members for inclusion into the final project. Difficulties understanding the theoretical questions were discussed along with a plan for tacking those problems at the next meeting. Drafts of the new report were circulated through email and improved in preparation for the last meeting.

At the sixth and final meeting everyone in the team read the report that had been polished over email. One error was found and fixed. All team members then agreed that the report was finished and ready to turn in.

The report you are reading here was written represents the independent work of the team members listed below and, in particular, does not contain the work of any other student in the class.

**2.** Given  $lpha \in [1,2]$  consider the recurrence relation

$$x_{n+1}=\left(1+rac{1}{x_n}
ight)^n$$

where  $x_1=lpha$ . Set lpha=1 and write a computer program in Julia that computes  $x_n$  for  $n=1,\ldots,50$ . Please include the full program listing and output. For reference,

$$x_{15} \approx 1.1.2634346914781789$$

and

 $x_{16} \approx 6283.87526674291.$ 

```
In [1]:
         alpha=1
Out[1]: 1
In [2]:
         x=alpha
         for n=1:49
             x=(1+1/x)^n
             println(n+1," ",x)
         end
        2 2.0
        3 2.25
        4 3.0137174211248285
        5 3.1461328651361042
        6 3.9749418398300125
        7 3.8436459461081727
        8 5.046648974589512
        9 4.24710770652293
        10 6.705641594809613
        11 4.014991943105442
        12 11.546255851436948
        13 2.7094151558091024
        14 59.372935430968965
        15 1.2634346914781789
```

- 16 6283.87526674291
- 17 1.0025492407161776
- 18 128267.71346883844
- 19 1.0001403407948934
- 20 523589.5387780858
- 21 1.0000381985532598
- 22 2.0963110569382557e6
- 23 1.000010494678316
- 24 8.387595658020036e6
- 25 1.0000028613721854
- 26 3.355323187755784e7
- 27 1.0000007748883877
- 28 1.3421632395741531e8
- 29 1.00000020861846
- 30 5.3686928798555815e8
- 31 1.0000000558795226
- 32 2.1474817879952881e9
- 33 1.0000000149011754
- 34 8.589932479998269e9
- 35 1.000000003958121
- 36 3.435973598800008e10
- 37 1.000000001047738
- 38 1.3743895080800003e11
- 39 1.0000000002764864
- 40 5.49755810924e11
- 41 1.000000000727596
- 42 2.199023252272e12
- 43 1.000000000190994
- 44 8.796093018596e12
- 45 1.0000000000050022
- 46 3.5184372084872e13
- 47 1.0000000000013074
- 48 1.40737488351004e14
- 49 1.000000000000341

Note that the output agrees with the expected values for  $x_{15}$  and  $x_{16}$ .

Based on the numerical evidence in the previous step, make a conjecture regarding the value of the limits

$$\lim_{n o\infty}x_{2n}$$

and

$$\lim_{n o\infty}x_{2n+1}$$

when  $x_1=1$ . Explain your reasoning in as much mathematical detail as possible.

A visual inspection of the output from the program in question 2 indicates that the even terms growing without bound while the odd terms appear to decrease towards 1.

Note that if  $x_{2n+1}=1+\epsilon$  for  $\epsilon$  very small then

$$1+\frac{1}{x_{2n+1}}=1+\frac{1}{1+\epsilon}=\frac{2+\epsilon}{1+\epsilon}\approx 2$$

Therefore

$$x_{2n+2} = \left(1 + rac{1}{x_{2n+1}}
ight)^{2n+1} pprox 2^{2n+1} o \infty$$

as  $n o \infty$  .

On the other hand if  $x_{2n}$  is very close to infinity than

$$1 + \frac{1}{x_{2n}} \approx 1$$

and so it's reasonable (but not guaranteed because  $1^\infty$  is an indeterminate form) that

$$x_{2n+1} = \left(1 + \frac{1}{x_{2n}}\right)^{2n} \approx 1^{2n} = 1$$

as  $n o \infty$  .

At anyrate, given the numerical evidence and the above heuristic argument, our conjecture is that

$$\lim_{n o\infty}x_{2n}=\infty$$

and

$$\lim_{n\to\infty} x_{2n+1} = 1$$

when  $x_1=1$ .

**4.** Change the program in the previous step to compute the values of  $x_n$  when  $\alpha=2$ . Look at the output and now make a conjecture regarding the value of the limits

$$\lim_{n o\infty}x_{2n}$$

and

$$\lim_{n o\infty}x_{2n+1}$$

when  $x_1=2$ . Explain the reasoning behind your conjecture.

```
In [3]:
         alpha=2
Out[3]: 2
In [4]:
         x=alpha
         for n=1:49
             x=(1+1/x)^n
             println(n+1," ",x)
         end
        2
          1.5
        3
           2.77777777777777
           2.515456000000001
        5
           3.814694524582706
        6
           3.202910914929112
        7
           5.105425602411104
           3.4977433630005117
           7.475734025618049
            3.0953609877250656
            16.436481006193045
        12
            1.9149171433392773
        13
            154.77995881694767
        14
            1.0873244478087112
        15
            9230.628993458884
        16
            1.001626257595803
        17
            64689.92026265085
        18
            1.0002628246012701
        19
            261524.77368382475
        20
            1.0000726533613016
            1.0478144924485891e6
        21
        22
            1.0000200419068173
        23
            4.193379435467222e6
            1.0000054848507496
        25
            1.6776111794572083e7
            1.0000014902152625
        27
            6.710756392755206e7
        28
            1.000000402339224
        29
            2.6843394397513205e8
        30
            1.000000108034036
        31
            1.0737400839916067e9
        32
            1.0000000288710484
        33
            4.2949653119971633e9
        34
            1.0000000076834112
        35
            1.7179866940000143e10
            1.0000000020372681
        37
            6.8719474216000046e10
        38
            1.0000000005384209
        39
            2.74877904132e11
```

- 40 1.000000001418812
- 41 1.099511624656e12
- 42 1.0000000000372893
- 43 4.39804650766e12
- 44 1.000000000009777
- 45 1.7592186040632e13
- 46 1.000000000002558
- 47 7.0368744173524e13
- 48 1.000000000000668
- 49 2.81474976706144e14
- EQ 1 000000000000174

When  $\alpha=2$  the subsequence which remains bounded and the one which goes to infinity appears to change. Likely, since 2 is bigger, then starting  $x_1$  with that larger number now makes the odd sequence  $x_{2n+1}$  be the one which tends to infinity while the evens sequence tends to 1. In particular, based on the numerical evidence, we conjecture that

$$\lim_{n o\infty}x_{2n}=1$$

and

$$\lim_{n o\infty}x_{2n+1}=\infty.$$

**5.** [Extra Credit] Use rigorous mathematical analysis to prove the conjectures made in the previous two steps.

The theory in the reference

J. Ewing, C Foias, An Interesting Serendipitous Real Number, Finite versus Infinite:
 Contributions to an Eternal Dilema, Springer-Verlag, pp. 119-120, 2000.

implies that limits conjectured in the previous questions are correct. Truthfully, more should be written for extra credit, for example quote the theorem from that book, to summarize the result and explain how the above conjectures follow from it. For this solution key, I have not gone into these necessary details.

More information may be found at

- Folas Constant Wikipedia, dowloaded October 17, 2020. https://en.wikipedia.org/ /wiki/Folas\_constant.
- 6. Define

$$lpha_* = \sup \Big\{ lpha : |x_{2n+1}| ext{ is bounded as } n o \infty \Big\}.$$

Intuitively,  $\alpha_*$  is the largest value of  $\alpha$  such that  $|x_{2n+1}|$  is bounded. Explain in details how the interval bisection method could be used to approximate  $\alpha_*$ .

The interval bisection method is a natural search technique that can be used to find a critical value at which a function or system or changes state or behavior. A typical example of such a change is when a continuous function f changes from being positive to negative. Note to get started that it is important to have two reference values a and b such that f(a) and f(b) have opposite signs. In this case, finding the critical value at which the transition occurs is equivalent to finding a root x such that f(x) = 0.

In the present case the value lpha=1 leads to the odd terms  $x_{2n+1}$  being bounded while lpha=2 leads to that sequence not being bounded. Bisection can be used to find the critical value at which the transition from bounded to unbounded behavior occurs. Thus, we take a=1 and b=1 and note that  $lpha_*\in [a,b]$ 

To search for  $\alpha_*$  we narrow down that value by bisecting the interval into the pieces [a,c] and [c,b] where c is given as the midpoint c=(a+b)/2. Now, if  $\alpha=c$  results in bounded behavior of the sequence, then we know  $\alpha_*\in [c,b]$ . On the other hand, if  $\alpha=c$  results in unbounded behavior of the sequence, then we know  $\alpha_*\in [a,c]$ .

In either case, the above bisection results in better knowledge of the true value of  $\alpha_*$ . The bisection algorithm can then be used to divide subsequence intervals in half until a bound that determines  $\alpha$  to the desired precision is obtained.

7. Write a computer program that bisects the interval [1,2] to find an approximation of  $lpha_*$  good to at least 4 significant digits. Include the full program listing and output.

First create a function that returns true if the resulting sequence is bounded and false if it isn't. This function assumes the behavior observed in the previous part of this assignment is typical and simply checks whether  $|x_{51}|$  is greater than or less than 2 in order to determine whether the entire sequence  $x_{2n+1}$  is bounded or not.

```
In [5]:
    function isoddbounded(alpha)
        x=alpha
        for n=1:50
              x=(1+1/x)^n
        end
        return abs(x)<2
        end</pre>
```

## Out[5]: isoddbounded (generic function with 1 method)

Before continuing, we then check that our function with the known values lpha=1 and lpha=2 to make sure it is working.

## Out[7]: false

Since  $lpha_*\in[1,2]$  we know the first digit of lpha is be one and the remaining digits will appear after the decimal point. Therefore, to obtain 4 significant digits we need to arrive at an interval [a,b] such that  $lpha_*\in[a,b]$  and with a length |b-a|<0.0005. Since the interval will be cut in half at each iteration, this implies at least n iterations where

$$\frac{1}{2^n} < 0.0005$$

will be sufficient. Taking logarithms as solving for  $m{n}$  yields

$$n > \log(1/0.0005)/\log(2)$$
.

Since

```
In [8]: log(1/0.0005)/log(2)
```

Out[8]: 10.965784284662087

we perform n=11 bisection steps.

```
In [9]:
    a=1
    b=2
    for n=1:11
        c=(a+b)/2
        if isoddbounded(c)
            a=c
        else
            b=c
        end
    end
    println("alpha_* in [",a,",",b,"]")
```

alpha\_\* in [1.18701171875,1.1875]

Therefore, the  $lpha_*pprox 1.187$  to 4 significant digits.

**8.** Explain theoretically how many times the interval [1,2] needs to be bisected to ensure the resulting approximation is good to at least 4 significant digits. What if 6 significant

digits are desired? How about 8 significant digits?

Most of the explanation is already contained in the answer to the previous question. In summary, at the beginning it is known that  $lpha_*\in[1,2]$  which is an interval of length 1 and further implies the first digit of  $lpha_*$  is one.

In general, since the first digit is one and the rest are after the decimal point, to obtain k significant digits it is required that the length of the interval be less than  $5 \times 10^{-k}$ . Consequently, taking the number of interations n such that

$$\frac{1}{2^n} < 5 \times 10^{-k}$$

or equivalently so

$$n > rac{\log(0.2 imes 10^k)}{\log 2} = 1 + (k-1)rac{\log 10}{\log 2}.$$

```
In [10]: Taking $k=4,6$ and $8$ yields
```

Out[10]: \$ (generic function with 1 method)

```
for k=4:2:8
    n=Integer(ceil(1+(k-1)*log(10)/log(2)));
    println("For $k significant digits use $n iterations")
end
```

```
For 4 significant digits use 11 iterations
For 6 significant digits use 18 iterations
For 8 significant digits use 25 iterations
```

9. Define

$$lpha^* = \inf \Big\{ lpha : |x_{2n}| ext{ is bounded as } n o \infty \Big\}.$$

Intuitively  $\alpha^*$  is the smallest value of  $\alpha$  such that  $|x_{2n}|$  is bounded. Modify the program from the previous step to obtain an approximation of  $\alpha^*$ . Include the output and describe what modifications were made to the code.

The main difference is that the function *isoddbounded* needs to be replaced by *isevenbounded* which stops at the 50th iteration rather than the 51st iteration.

```
In [12]:
    function isevenbounded(alpha)
        x=alpha
        for n=1:49
            x=(1+1/x)^n
        end
        return abs(x)<2
    end</pre>
```

Out[12]: isevenbounded (generic function with 1 method)

Then since

```
In [13]: isevenbounded(1)

Out[13]: false
In [14]: isevenbounded(2)
```

Out[14]: true

One also needs to switch the logic in the if statement to reassign the upper index b of the interval when  $x_{2n}$  is bounded and lower index a when it's not.

```
In [15]:
    a=1
    b=2
    for n=1:11
        c=(a+b)/2
        if isevenbounded(c)
            b=c
        else
            a=c
        end
    end
    println("alpha^* in [",a,",",b,"]")
```

alpha^\* in [1.18701171875,1.1875]

**10.** Based on the numerical evidence obtained in the previous two problems conjecturewhether the values of  $\alpha_*$  and  $\alpha^*$  are equal or different.

After 11 iterations  $\alpha_*$  are and  $\alpha^*$  are still in the same interval. This numerical evidence suggests that  $\alpha_*=\alpha^*$ . For greater confirmation, we check after 25 iterations.

```
In [16]:
    a=1
    b=2
    for n=1:25
        c=(a+b)/2
        if isoddbounded(c)
            a=c
        else
            b=c
        end
end
println("alpha_* in [",a,",",b,"]")
```

alpha\_\* in [1.187452346086502,1.1874523758888245]

```
In [17]:
    a=1
    b=2
    for n=1:25
        c=(a+b)/2
        if isevenbounded(c)
            b=c
        else
            a=c
        end
    end
    println("alpha^* in [",a,",",b,"]")
```

alpha^\* in [1.187452346086502,1.1874523758888245]

Even for this much smaller interval we numerically see that  $\alpha_*$  and  $\alpha^*$  are still in the same interval. This further suggests they are equal.

11. [Extra Credit] Use rigorous mathematical analysis to prove the conjecture stated in the previous step. Alternatively, further support your conjecture by computing additional digits of  $\alpha_*$  and  $\alpha^*$  using the BigFloat arbitrary precision arithmetic built into the Julia programming language.

The reference cited in question 5 proves that  $\alpha_*=\alpha^*$ . For extra credit, an more explanation how the theory in those works implies  $\alpha_*=\alpha^*$  and a brief description of the techniques used, if not a rephrasing of the full proof. An original proof of this result would be even better.

We finish with an example of using BigFloat arbitrary precision arithmetic to show that  $\alpha_*$  and  $\alpha^*$  agree to many many digits.

```
In [18]:
          function bigoddbounded(alpha)
              x=alpha
               for n=1:250
                   x=(1+1/x)^n
               end
               return abs(x) < 2
          end
          function bigevenbounded(alpha)
              x=alpha
               for n=1:249
                   x=(1+1/x)^n
               end
               return abs(x) < 2
          end;
In [19]:
          a=big(1)
          b=big(2)
          for n=1:256
              c=(a+b)/2
               if bigoddbounded(c)
                   a=c
               else
                   b=c
               end
          end
          println("alpha * in [\n",a,",\n",b,"]")
         alpha * in [
         1.187452351126501054595480158396519351215692681585860353010104126198780
         418723321,
         1.187452351126501054595480158396519351215692681585860353010104126198780
         418723338]
In [20]:
          a=big(1)
          b=big(2)
          for n=1:256
              c = (a+b)/2
               if bigevenbounded(c)
                   b=c
               else
                   a=c
               end
          end
          println("alpha^* in [\n",a,",\n",b,"]")
         alpha^* in [
         1.187452351126501054595480158396519351215692681585860353010104126198780
         418723321,
```

	The intervals are the same after 256 interations. This implies that $lpha_*=lpha^*$ agree to within
In [21]:	b-a
Out[21]:	1.727233711018888925077270372560079914223200072887256277004740694033718 360632485e-77
	or equivalently are the same to at least $77$ decimal digits.
Tn [ ].	

1.187452351126501054595480158396519351215692681585860353010104126198780